

BISS-E

— Basic Interoperable Scrambling System
with Encrypted keys

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Summary

This specification describes a scrambling system for use on digital contribution circuits (satellite, DSNG etc.) which use MPEG-2 compression and the DVB-S modulation scheme.

The use of encrypted session keys provides a high degree of security and allows the implementation of a centrally-managed conditional-access system which is suitable for *Eurovision* applications.

1. Introduction

1.1. Overview

The rapid increase in the use of Digital SNG (DSNG) technology has resulted in the offering of digital codec equipment by a number of vendors. At the same time, the absence of standard methods for securing and scrambling DSNG broadcasts has spawned the development of several different proprietary security mechanisms.

The widespread acceptance of DVB standards makes possible the proposal and provision of a security mechanism which offers interoperability between the equipment of different DSNG vendors. This would enable broadcasters to combine equipment from among several vendors, while making systems more future-proof.

In June 1999, the DVB Crypto Experts Group agreed on a standard mechanism for the scrambling of DSNG transmissions, based on the DVB Common Scrambling Algorithm (CSA). The Basic Interoperable Scrambling System (BISS) specification is based on a modified DVB-CSA and the use of keys called Session Words (see Tech doc. 3290).

“BISS with Encrypted keys” (BISS-E) introduces an additional mechanism that equipment shall implement (in software) to accept the insertion of Encrypted Session Words while conserving their interoperability. This mechanism is backward compatible with the BISS specification, mode 1.

For the sake of completeness, this document is self-contained in that it includes the complete specification for BISS *mode 1* and the new *mode E*.

1.2. Nomenclature

Throughout this document, the following definitions are used:

- * **Scrambler** relates to the overall mechanisms required to meet the DVB Common Scrambling Specification (Part 2).
- * **Scrambling Module** relates to the Super Scrambling Mechanisms required to meet the DVB Common Scrambling Specification (Part 2).
- * **Session Key** relates to the key that is unique and constant for the duration of the transmission. This may be a fixed CW, used for scrambling the transport stream directly or for adding a level of indirection – a key which is used to scramble changing CWs within Entitlement Control Messages.
- * **Session Word** relates to the word from which the Session Key is derived, i.e. the Session Word is not used directly in the scrambling process, but is transformed by some mechanism into the Session Key.
- * **Unit** relates to a device for which this specification might apply.
- * **Managing Centre** refers to an organisation controlling or managing the conditional access system.
- * **Decryption function** refers to a logical function used to decrypt Encrypted Session Words with the help of a key. It produces a 48-bit result from which the clear Session Word is derived.
- * **Interoperable function** refers to a decryption function that shall be embedded in all units.

1.3. Notations

The bits in binary numbers or sequences are numbered from the left, according to the engineering notation. Bit 0 is on the right and is the least significant one; the left bit is the most significant bit.

Example for an n-bit number:

$$b_{n-1}b_{n-2} \dots b_1b_0$$

Sequences of bytes are numbered starting at zero for the first byte, which is the MSB.

1.4. Security Requirements

The DSNG model requires the direct entry of a Session Word at the transmitter and receiver, to control access to the transmission. The sender and receiver(s) of the transmission share the Session Word, such that only the intended parties will receive the transmission, outlined as follows:

- 1) Session Word entered at the DSNG unit in the field or at the transmitting earthstation.
- 2) Session Word entered at the receiving IRDs.
- 3) If the Session Words are the same, then the IRDs are able to decrypt the broadcast.
- 4) If the Session Words are different, the broadcast is not received.

The security requirements for fixed contribution systems are somewhat different to the DSNG model. The secure exchange of Session Keys is fundamental to such systems and is achievable. For fixed systems requiring interoperability with DSNG units, external control systems may be employed to allow the transmission of Entitlement Management Messages (EMMs) for securely exchanging Session Keys between transmitting and receiving sites. This model works for transmission sites that are part of the fixed network, but when receive sites are accepting a transmission from a DSNG unit, the operation must revert to the direct-entry method described above.

1.5. Modes of Operation

The Scrambler must be capable of supporting the following four modes of operation:

- * **Mode 0:** No scrambling.
- * **Mode 1:** All components are scrambled by a fixed control word (CW), derived from a clear Session Word (SW).
- * **Mode E:** All components are scrambled by a fixed CW, derived from an Encrypted Session Word (ESW).

The Scrambler shall implement the Super Scrambling operations as defined in the DVB Common Scrambling Specification (Part 2). The scrambling mechanism shall be applied at Transport level only.

To support the various modes of operation, the Scrambler must be capable of inserting ECM streams into the multiplex and these streams shall be appropriately identified within the PMT. The use of EMM streams has no application within the modes of operation described within this document.

A CAT shall be present in the multiplex for BISS mode 1 and BISS-E, although the table shall be empty, as no EMM stream will be present.

A Scrambler that only supports a subset of the defined modes of operation must do so according to an imposed hierarchy. A Scrambler providing support for mode E must also support modes 0 and 1. Likewise, a Scrambler providing support for mode 3, must also support modes 0, 1 and 2.

1.6. Mode 0

The scrambler must be capable of disabling the scrambling operation. In this mode there will be no *CA_descriptor* in the PMT and no ECM stream. The *Transport_Scrambling_Control* bits of the Transport Packets will be set to “00”.

2. BISS Mode 1 — functional requirements

2.1. Overview

In this mode, the Scrambler uses a fixed Control Word (CW) for the duration of the transmission. The operator shall enter a Session Word, which is transformed into the Session Key (SK) for use by the Scrambling Module. In this mode, the terms Session Word and Session Key are synonymous with the terms Control Word and Common Key from the DVB Common Scrambling specifications, respectively. An overview is given in *Fig. 1*.

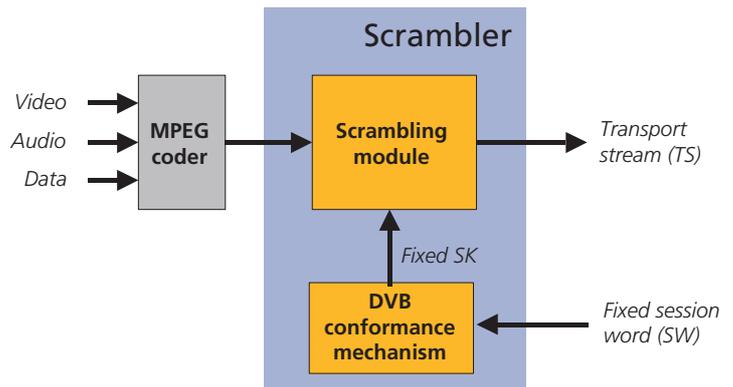


Figure 1
Overview of BISS Mode 1.

The SW is a 48-bit word which is transformed by the Scrambler into a 64-bit SK using the Conformance Mechanism defined as part of the DVB Common Scrambling specifications.

The 48-bit SW is first mapped to the 64-bit CW by the Scrambler, prior to applying the Conformance Mechanism. The mapping of bytes between the 48-bit SW and the 64-bit CW is given in *Table 1*.

Table 1
SW to fixed CW mapping.

64-bit CW	48-bit SW
CW(1)	SW(1)
CW(2)	SW(2)
CW(3)	SW(3)
CW(4)	see note ^a
CW(5)	SW(4)
CW(6)	SW(5)
CW(7)	SW(6)
CW(8)	see note ^b

- a. CW(4) is derived from SW(1)..SW(6) by the DVB-defined Conformance Mechanism.
- b. CW(8) is derived from SW(1)..SW(6) by the DVB-defined Conformance Mechanism.

In this mode there will be a *CA_descriptor* in the PMT, present at programme level, but no ECM stream. A single unique *CA_System_ID* is assigned to identify mode 1.

The *Transport_Scrambling_Control* bits of the Transport Packets shall be set to “10”.

Manual entry of the SW shall be in Hexadecimal, with the digits entered most-significant-nibble first, i.e. from left to right as viewed in hexadecimal notation.

For example, 0xA13DBC42908F, would be entered in the following sequence: A,1,3,D,B,C,4,2,9,0,8,F.

Remote entry of the SW shall also be provided, although the specification of this interface is beyond the scope of this document.

The Scrambler shall ensure that the SK used by the Scrambling Module cannot be changed more than 10 times in a 5 minute period and that there is a minimum of 10 seconds between changes.

2.2. CA_descriptor

The *CA_descriptor* which must be present in the PMT to support mode 1 is defined in Table 2. Semantics:

CA_system_ID: this is a 16-bit field indicating the type of CA system applicable for the associated ECM streams. The value of this field for mode 1 is 0x2600.

CA_PID: this is a 13-bit field indicating the PID of the Transport Stream packets that shall contain the ECM information. For mode 1, no ECM information is required, so this field shall contain the value 0x1FFF.

Table 2
Conditional access descriptor – Mode 1.

Syntax	No. of bits	Identifier
CA_descriptor() {		
descriptor_tag	8	uimsbf
descriptor_length	8	uimsbf
CA_system_ID	16	uimsbf
reserved	3	bslbf
CA_PID	13	uimsbf
}		

3. BISS Mode E — functional requirements

3.1. Clear Session Word

The unit shall be BISS-Mode 1-compliant. It shall support the insertion of clear 48-bit Session Word (SW) through the front panel and through the remote control interface. It shall use the SW as specified in Chapter 2 of this document.

The clear session word, once entered via the user interface or remote control port, shall not be readable through any unit interface.

3.2. Encrypted Session Word (ESW)

The unit shall support the insertion of Encrypted Session Words through the front panel and through a remote control interface. The definition of the remote control port is outside the scope of this recommendation.

The Encrypted Session Word is a 64-bit number that is transformed by the unit into a 48 bit clear Session Word. The clear Session Word is then used by the unit to decrypt the broadcast according to Chapter 2 (BISS Mode 1).

Once the Encrypted Session Word has been entered via the front panel or via the remote control interface, it shall be impossible to read it back through any unit interface.

The manual entry of the ESW shall be in hexadecimal form, with the 16 digits entered most-significant nibble (i.e. the leftmost nibble) first.

For example, if the ESW is 0xF76EE249BE01A286, it shall be entered in the following sequence:

F, 7, 6, E, E, 2, 4, 9, B, E, 0, 1, A, 2, 8 and 6.

3.3. Decryption Scheme

3.3.1. Overview

The equipment shall include the following features:

- * An identifier, denoted ID, comprising a 56 bit hexadecimal word which shall be injected by the user and shall be used as the default. The injected ID is mandatory. Optionally, in addition, the supplier may bury an ID. In this case, the user shall actively select the buried ID.
- * A DES decryption function denoted $f()$, as described in Section 3.3.3. Additional functions may be supplied but are beyond the scope of this document.
- * A simple post-processing function denoted $P()$, as described in Section 3.3.4.

An illustration of the processing of the ESW in the unit to provide the clear Session Word is given in Figure 1 and further examples are described in the Appendix. The role of Management Centre generating the ESW encrypted by DES is indicated but the detailed specification is outside of the scope of this document.

The mapping of the ID is a simple expansion from 56 to 64 bits by adding an odd parity bit after every 7 bits. The reduction of the decrypted session word for 64 to 48 bits is obtained by deleting the first and last bit of each byte (see examples in the Appendix).

After the application of the post-processing function $P()$, the clear session word SW is obtained to feed the BISS equipment as in Mode 1.

3.3.2. Unit identifiers

This document specifies two types of identifiers for each unit.

- 1) An **injected identifier** ID_i that is a secret key embedded in the unit. This is mandatory.
- 2) Additionally, the manufacturer may provide a **buried identifier** ID_b , defined by the manufacturer and uniquely linked to the device itself. This is not mandatory, but if implemented it shall comply with this document.

A user shall be able to select the identifier of his choice via the front panel and the remote control interface. The selected identifier is used as the Active key to decrypt the Encrypted Session Word.

a) Injected ID – Mandatory

The injected ID is a 56-bit identifier that can be entered by the user at any time in the BISS unit.

Units shall support the insertion of the injected ID through its front panel and through its remote control interface. There shall be no mechanism for reading back part or all of the injected ID via any unit interface.

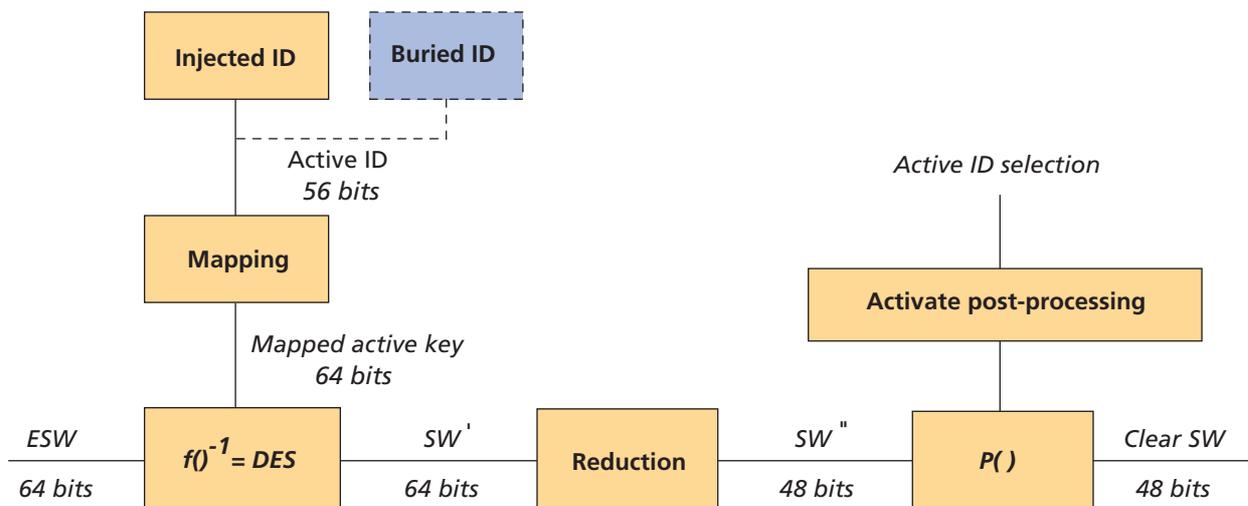


Figure 2
The signal processing required to produce a clear Session Word.

The same ID can be injected in more than one piece of equipment, e.g. for redundancy management.

The manual or remote entry of the injected ID shall be in hexadecimal form, and the 14 digits are entered with the most-significant nibble (i.e. the left-most nibble) first.

For example, if the injected ID is 0xF09A423F56738A, it shall be entered in the following sequence:

F, 0, 9, A, 4, 2, 3, F, 5, 6, 7, 3, 8 and A.

b) Buried ID – Option for the supplier

The buried ID is a 56-bit identifier that uniquely identifies a particular unit. Buried IDs are optional.

Two different units shall have a different buried ID, at least for the equipment produced by the same manufacturer. Units from different manufacturers could have the same buried ID, but it shall be a fortuitous case.

The manufacturer shall ensure that without his agreement, nobody can modify the buried ID.

3.3.3. Decryption function

Units shall implement the interoperable function specified below. Additional decryption functions may optionally be implemented. In this case, it shall be possible to select the decryption function via the front panel and the remote control interface. The definition of the additional functions is outside the scope of this document.

The interoperable function is mandatory on all units. It uses the 64-bit ESW and the 64-bit mapped active key to compute a 48-bit word called SW''.

a) Mapping

The mapped active key is derived from the 56-bit active ID by mapping the 56 active ID bits into the upper 7 bits of a sequence of 8 bytes, then the LSB of each byte is set such that the resultant byte has odd parity.

**Table 3
Active ID to mapped Active Key mapping.**

Active Key Bit	Mapped Byte / Bit	Mapped Key Bit
55	0 / 7	63
54	0 / 6	62
53	0 / 5	61
52	0 / 4	60
51	0 / 3	59
50	0 / 2	58
49	0 / 1	57
Odd_parity	0 / 0	56
48	1 / 7	55
Etc.	Etc.	Etc.

[Note that the above uses engineering notation: msb = bit 63, lsb = bit 0]

b) Interoperable function

The interoperable function is the simple DES algorithm used in ECB mode and decrypt state. This function is described in the FIPS PUB 46-3 [2] and FIPS PUB 81 [3]. The algorithm key is the 64-bit mapped active key ¹. The data bloc to decrypt is the ESW. Note that this document uses the engineering notation (i.e. msb = bit 63, lsb = bit 0), while DES uses FIPS notation (i.e. msb = bit 1, lsb = bit 64). Hence the bit mapping is:

- * DES key (1..64) ≤ mapped active key (63..0)
- * DES data bloc (1..64) ≤ ESW (63..0)

The result of the decryption algorithm is defined on 64 bits and is called **SW''**.

The mapping between **SW'** and **SW''** is given in *Table 4*. This mapping simply removes the most significant and the least-significant bit of each byte. The third column of the table is the DES result with the bits numbered with FIPS notation.

Table 4
Mapping between SW' and SW''

SW''(47)	SW'(62)	D(2)
SW''(46)	SW'(61)	D(3)
SW''(45)	SW'(60)	D(4)
SW''(44)	SW'(59)	D(5)
SW''(43)	SW'(58)	D(6)
SW''(42)	SW'(57)	D(7)
SW''(41)	SW'(54)	D(10)
SW''(40)	SW'(53)	D(11)
SW''(39)	SW'(52)	D(12)
SW''(38)	SW'(51)	D(13)
SW''(37)	SW'(50)	D(14)
SW''(36)	SW'(49)	D(15)
SW''(35)	SW'(46)	D(18)
SW''(34)	SW'(45)	D(19)
SW''(33)	SW'(44)	D(20)
SW''(32)	SW'(43)	D(21)
SW''(31)	SW'(42)	D(22)
SW''(30)	SW'(41)	D(23)
SW''(29)	SW'(38)	D(26)
SW''(28)	SW'(37)	D(27)
SW''(27)	SW'(36)	D(28)

1. The DES algorithm does not use bits 8, 16, 24, 32, 40, 48, 56 and 64 of the key, numbered with the FIPS notation (see [2]). Hence the useful key length is actually of 56 bits, which is compatible with exportation restrictions.

Table 4
Mapping between SW' and SW''

SW''(26)	SW'(35)	D(29)
SW''(25)	SW'(34)	D(30)
SW''(24)	SW'(33)	D(31)
SW''(23)	SW'(30)	D(34)
SW''(22)	SW'(29)	D(35)
SW''(21)	SW'(28)	D(36)
SW''(20)	SW'(27)	D(37)
SW''(19)	SW'(26)	D(38)
SW''(18)	SW'(25)	D(39)
SW''(17)	SW'(22)	D(42)
SW''(16)	SW'(21)	D(43)
SW''(15)	SW'(20)	D(44)
SW''(14)	SW'(19)	D(45)
SW''(13)	SW'(18)	D(46)
SW''(12)	SW'(17)	D(47)
SW''(11)	SW'(14)	D(50)
SW''(10)	SW'(13)	D(51)
SW''(9)	SW'(12)	D(52)
SW''(8)	SW'(11)	D(53)
SW''(7)	SW'(10)	D(54)
SW''(6)	SW'(9)	D(55)
SW''(5)	SW'(6)	D(58)
SW''(4)	SW'(5)	D(59)
SW''(3)	SW'(4)	D(60)
SW''(2)	SW'(3)	D(61)
SW''(1)	SW'(2)	D(62)
SW''(0)	SW'(1)	D(63)

3.3.4. Post-processing function

The post-processing function $P()$ converts SW'' into the clear SW. The conversion gives different results depending on the type of the ID used to decrypt the ESW.

When the active key is the injected ID, function $P()$ is the identity function (i.e. SW'' = SW).

If the optional ; buried ID is used, function $P()$ consists of rotating the SW'' of one bit to the right. In the other cases, the definition of $P()$ is outside the scope of this document. It shall however have a different mathematical behaviour (i.e. it shall produce different results) than the expressions used when the active key is the injected ID or the buried ID.

If $SW'' = b_{47} b_{46} \dots b_1 b_0$

$SW = P(SW'') =$

$b_{47} b_{46} \dots b_1 b_0$	if the active key is an injected ID
$b_0 b_{47} b_{46} \dots b_2 b_1$	if the active key is a buried ID
undefined	otherwise

Bibliography

- [1] Recommendation H.222.0, ISO/IEC 13818-1: **Information Technology – Generic coding of moving pictures and associated audio: Systems.**
- [2] ETR 162: **Digital broadcasting systems for television, sound and data services; Allocation of Service Information (SI) codes for Digital Video Broadcasting (DVB) systems.**
- [3] ETR 289: **Digital broadcasting systems for television, sound and data services; Support for use of scrambling and conditional access (CA) within digital broadcasting systems.**
- [4] ETS 300 468: **Digital broadcasting systems for television, sound and data services; Specification for Service Information (SI) in Digital Video Broadcasting (DVB) systems.**
- [5] FIPS PUB 46-2: **Data Encryption Standard.**
- [6] FIPS PUB 46-3: **Data Encryption Standard.**
- [7] FIPS PUB 81: **DES Modes of Operation.**
- [8] <http://www.etsi.org/>: **DVB Common Scrambling Specifications.**

Abbreviations

bslbf	Bit String, Left Bit First	KE	Key Escrow
CA	Conditional Access	lsb	Least Significant Bit
CAT	Conditional Access Table	LSB	Least Significant Byte
CK	Common Key	MC	Management Centre
CW	Control Word	msb	Most Significant Bit
DES	Data Encryption Standard	MSB	Most Significant Byte
DSNG	Digital SNG	PMT	Programme Map Table
ECB	Electronic Codebook	SK	Session Key
ECM	Entitlement Control Message	SNG	Satellite News Gathering
EDE	Encode, Decode, Encode	SW	Session Word
EMM	Entitlement Management Message	uimsbf	Unsigned Integer, Most Significant Bit First
ESW	Encrypted Session Word		

Appendix A DES usage in BISS-E

The following tables show six examples to the processes described later on in 2 cases:

“Active ID” is “injected ID” (table A.1)

“Active ID” is “buried ID” (table A.2)

Table A.1 — Results while using “injected ID”.

Encrypted Session Word (64 bits) (Note 1)	Active ID (56 bits) (Note 2)	Expanded Session Word (64 bits) (Note 4)
	Mapped Active Key (64 bits) (Note 3)	Session Word (after reduction) (48 bits) (Note 5)

Example 1:

E81816B87E5CF9C4	FB5F9C585DD359	B2CB8F3948770EF9
	FBAEE68A85EF4CB3	6651DC93B1FC

Example 2:

FE0E810E96648C43	DC91106F13C87B	AA8130F0B47D381B
	DC49450DF19E20F7	5406386BE70D

Example 3:

51B276C10BA89683	89F27FB3ACA107	066283A18293C5A6
	89F89EF73B64850E	0F1050049893

Example 4:

1D599C25C30C2C11	85513F83A5FF36	0A74BCCBCA80AAFE
	85A84FF13B2FFD6D	17A7A594057F

Example 5:

ED89E0D818B43B98	CD23BF7D8F6BC3	840D91C0F6444C27
	CD91EFefd97AAE86	086220EE2993

Example 6:

DA457923D39BDB81	E0A48CAA093AC7	7EE18611180D64A2
	E0522394A149EA8F	FF00C8306C91

Table A.2 — Results while using “buried ID”.

Encrypted Session Word (64 bits) (Note 1)	Active ID (56 bits) (Note 2)	Expanded Session Word (64 bits) (Note 4)
		Session Word (after reduction) (48 bits) (Note 5)
	Mapped Active Key (64 bits) (Note 3)	Session Word (after post processing P) (48 bits) (Note 6)

Example 1:

389C30A2AB29F3F2	FB5F9C585DD359	E6149C72936D9F71
		CCA3B92763F8
	FBAEE68A85EF4CB3	6651DC93B1FC

Example 2:

288871F187278E6A	DC91106F13C87B	D500E2E1EBF8F135
		A80C70D7CE1A
	DC49450DF19E20F7	5406386BE70D

Example 3:

53C0C292CC2A88B3	89F27FB3ACA107	0E4405410427894C
		1E20A0093126
	89F89EF73B64850E	0F1050049893

Example 4:

6EC275343C161742	85513F83A5FF36	6EC275343C161742
		2F4F4B280AFE
	85A84FF13B2FFD6D	17A7A594057F

Example 5:

0E027A2DBDB5622C	CD23BF7D8F6BC3	8899A203EE8A194D
		10C441DC5326
	CD91EFefd97AAE86	086220EE2993

Example 6:

C5E9F325179F97A2	E0A48CAA093AC7	7FC10CA0311A4846
		FE019060D923
	E0522394A149EA8F	FF00C8306C91

Notes:

- 1) **ESW** is received from the Management Centre (64 bits) and is entered manually or via the remote control port.
- 2) **Active ID** is used for the DES⁻¹ decryption (56 bits).
- 3) The **Active ID** is expanded to 64 bits by taking each group of 7 bits, adding an odd parity bit and forming a byte (the parity bit is the least significant bit of the byte). This **Mapped Active Key** is used as the key of the DES⁻¹ decryption (in ECB mode).
- 4) **Expanded Session Word** is the result of the **DES⁻¹** process (64 bits).
- 5) **Session Word after reduction (SW'')** is calculated by taking each byte of **Expanded Session Word**, removing both its lsb and msb, and concatenating 8 times the 6-bit words to 48 bits.
- 6) According to the **Active ID Selection** (Buried or Injected), a post-processing function **P()** is activated on the **SW''**. The function **P()** is defined to be either **Identity** function (for Injected ID), or **rotate right by 1 bit SW''** (for the buried ID). The output of **P()** is the Session Word (48 bits).

The **Equipment** processes:

Getting the **Encrypted Session Word (ESW)** into the equipment:

- 1) Map the **Active ID''** to get the **Mapped Active Key**.
- 2) Do the decryption using DES⁻¹ (at ECB mode) with the **Mapped Active Key** as key, and **ESW** as data.
- 3) The result is the **Expanded Session word (SW')**.
- 4) Reduce the result to a 48-bit word as in *Note 5* above.
- 5) Do the post processing function as described by Note 6. The output of **P()** is denoted as the “**Session Word**” (**SW**).

The **Management Centre** processes:

Choosing a **Session Word**:

- 1) Do the preprocessing function (inverse of the post-processing function, described by Note 6), and get **SW''**.
- 2) Do the Expansion (expand each 6 bits to one byte, filling both the least and the most significant bits of each byte by 2 random bits) and get **Expanded Session Word (SW')**.
- 3) Map the **Active ID** to get the **Mapped Active Key**.
- 4) Do the encryption using DES (at ECB mode) with the **Mapped Active Key** as key, and **Expanded SW** as data.
- 5) Get the results of the encryption (**ESW**), and send it to the equipment.